- 1. Use the binary search algorithm to try to locate the name NIGEL in the following alphabetical list. Clearly indicate how you chose your pivots and which part of the list is being rejected at each stage.
  - 1. Bhavika
  - 2. Clive
  - 3. Elizabeth
  - 4. John
  - 5. Mark
  - 6. Nicky
  - 7. Preety
  - 8. Steve
  - 9. Trevor
  - 10. Verity

(Total 4 marks)

2. 6	550 ·	431	245	643	455	134	710	234	162	452
------	-------	-----	-----	-----	-----	-----	-----	-----	-----	-----

(a) The list of numbers above is to be sorted into **descending** order. Perform a Quick Sort to obtain the sorted list, giving the state of the list after each pass, indicating the pivot elements.

(5)

The numbers in the list represent the lengths, in mm, of some pieces of wood. The wood is sold in one metre lengths.

(b) Use the first-fit decreasing bin packing algorithm to determine how these pieces could be cut from the minimum number of one metre lengths. (You should ignore wastage due to cutting.)

(4)

(c) Determine whether your solution to part (b) is optimal. Give a reason for your answer.

(2) (Total 11 marks) 3.

Hajra	Vicky	Leisham	Alice	Nicky	June	Sharon	Tom	Paul
(H)	(V)	(L)	(A)	(N)	(J)	(S)	(T)	(P)

The table shows the names of nine people.

(a) Use a quick sort to produce the list of names in ascending alphabetical order.

You must make your pivots clear.

(4)

(b) Use the binary search algorithm on your list to locate the name Paul.

(4) (Total 8 marks)

4. A builder is asked to replace the guttering on a house. The lengths needed, in metres, are

Guttering is sold in 4 m lengths.

(a) Carry out a quick sort to produce a list of the lengths needed in **descending** order. You should show the result of each pass and identify your pivots clearly.

(5)

(b) Apply the first-fit decreasing bin-packing algorithm to your ordered list to determine the total number of 4 m lengths needed.

(4)

(c) Does the answer to part (b) use the minimum number of 4 m lengths? You must justify your answer.

(2) (Total 11 marks) 5. Miri Jessie Edward Katie Hegg Beth Louis Philip Natsuko Dylan (a) Use the quick sort algorithm to sort the above list into alphabetical order. (5) Use the binary search algorithm to locate the name Louis. (b) (4) (Total 9 marks) 6. Max Hannah Lauren John Kieran Tara Richard Imogen Use a quick sort to produce a list of these names in ascending alphabetical order. You (a) must make your pivots clear. (5) (b) Use the binary search algorithm on your list from part (a) to try to locate the name 'Hugo'. (4) (Total 9 marks)

- 7. Use the binary search algorithm to try to locate the name NIGEL in the following alphabetical list. Clearly indicate how you chose your pivots and which part of the list is being rejected at each stage.
  - 1. Bhavika
  - 2. Clive
  - 3. Elizabeth
  - 4. John
  - 5. Mark
  - 6. Nicky
  - 7. Preety
  - 8. Steve
  - 9. Trevor
     10. Verity

## (Total 4 marks)

8.		45,	56,	37,	79,	46,	18,	90,	81,	51			
	(a)	Using the numbers i	quick s into <b>asc</b>	sort algor <b>ending</b> c	ithm, p order.	erform <b>o</b>	ne comp	olete iter	ation tov	vards so	rting the	ese	(2)
		<b>TT T</b>		. 1	• •	G		1.		1	.1		(-)
	(b)	Using the list into <b>d</b>	bubble escendi	sort algo ing order	orithm,	perform	one con	iplete pa	ss towar	ds sortir	ig the oi	riginal	(2)
	Ano	ther list of r	numbers	s, in ascer	nding o	rder, is							
		7,	23,	31,	37,	41,	44,	50,	62,	71,	73,	94	
	(c)	Use the b	inary se	earch algo	orithm t	o locate	the num	ber 73 ir	n this list				
											(	Total 8 r	(4) narks)
9.													
				1		Glasgo	W						
				2		Newcas	stle						
				3		Manche	ster						
				4		York							
				5		Leicest	er						
				6	•	Birming	ham						
				7		Cardi	ff						
				8		Exete	r						
				9		Southam	pton						
				10	).	Plymou	ıth						
	A bi	nary search	is to be	e perform	ed on t	he names	s in the l	ist above	e to locat	te the na	me New	vcastle.	
	(a)	Explain w	vhy a bi	nary sear	ch canı	not be pe	rformed	with the	e list in i	ts preser	nt form.		(1)
	(b)	Using an	appropi	riate algo	rithm. a	alter the ]	list so th	at a bina	rv searc	h can be	perforn	ned.	
		State the	name of	f the algo	orithm y	ou use.			<b>)</b>		r		
													(4)
	(c)	Use the b	inary se	earch algo	orithm o	on your r	new list	to locate	the nam	e Newca	astle.		
													(4)

(4) (Total 9 marks) 10. The following list gives the names of some students who have represented Britain in the International Mathematics Olympiad.
Roper (R), Palmer (P), Boase (B), Young (Y), Thomas (T), Kenney (K), Morris (M), Halliwell (H), Wicker (W), Garesalingam (G).
(a) Use the quick sort algorithm to sort the names above into alphabetical order.
(5)
(b) Use the binary search algorithm to locate the name Kenney.

1. 
$$\left[\frac{1+10}{2}\right] = 6$$
 Nicky M1  
 $reject top of list$   
 $\left[\frac{7+10}{2}\right] = 9$  Trevor A1  
 $reject bottom of list$   
 $\left[\frac{7+8}{2}\right] = 8$  Steve A1  
 $reject bottom of list$   
 $[7] = 7$  Preety A1 4  
 $reject$ 

Nigel not in list

[4]

<b>2.</b> (a) E.g.
--------------------

<b>D</b> .5.				-	-				-	
650	431	245	643	455	710	234	162	452	134	M1
650	643	710	455	431	245	234	162	452	134	A1
650	710	643	455	431	245	452	234	162	134	A1 ft
710	650	643	455	431	452	245	234	162	134	A1 ft
710	650	643	455	452	431	245	234	162	134	A1
					•					

5

(b)	Bin 1 710 + 245	Bin 3 643 + 162 + 134	Bin 5 431 M1A1	
	Bin 2 650 + 234	Bin 4 455 + 452	A1A1(ft)	4

(c) $\frac{4116}{1000} = 4.116$	5 bins needed optimal	M1A1(ft)	2	
---------------------------------	-----------------------	----------	---	--

[11]

### **3.** (a)

Н	V	L	А	Ν	J	S	Т	Р	(N)	M1	
Н	L	А	J	N	V	S	Т	Р	(A, T)	A1	
А	Η	L	J	N	S	Р	Т	V	(L, P)	A1 ft	
А	Η	J	L	N	Р	S	Т	V	(J)		
Α	Η	J	L	Ν	Р	S	Т	V		A1 cso	4

# <u>Note</u>

1M1: quick sort, pivots, p, chosen and two sublists one p. 1A1: first pass correct and next pivots chosen correctly/consistently. 2A1ft: second pass correct, next pivots correctly/consistently chosen. 3A1: all correct, cso.

(b) 
$$1^{\text{st}} \text{ choice } \left[\frac{1+9}{2}\right] = 5 \text{ Nicky, reject } 1-5 \text{ M1 A1}$$
  
 $2^{\text{nd}} \text{ choice } \left[\frac{6+9}{2}\right] = [7,5] = 8 \text{ Tom reject } 8-9 \text{ A1}$ 

$$3^{rd} \text{ choice } \left[ \begin{array}{c} 2 \\ 2 \end{array} \right]^{-[7.5]=6} \text{ Four, reject } 6^{-5} \text{ for all results}$$

$$3^{rd} \text{ choice } \left[ \frac{6+7}{2} \right] = [6.5] = 7 \text{ Sharon, reject } 7$$

$$4^{th} \text{ choice } 6 \text{ Paul name found}$$

$$A1 \text{ cso} \quad 4$$

### <u>Note</u>

1M1: binary search on what they think is a alphabetical list, choosing pivot, rejecting half list.

1A1: first pass correct, condone 'sticky' pivot here, bod generous

2A1: second pass correct, pivot rejected.

3A1: cso.

Note: If incorrect list in (a) mark (b) as a misread.

### **Alternative solutions**

Middle right

Н	V	L	А	Ν	J	S	Т	Р	(N)	M1
Н	L	А	J	N	V	S	Т	Р	(A T)	A1
А	Н	L	J	N	S	Р	Т	V	(L P)	A1ft
А	Н	J	L	N	Р	S	Т	V	(J)	
А	Н	J	L	N	Р	S	Т	V	-	A1cso

list sorted

Midd	le left	Į								
Н	V	L	А	Ν	J	S	Т	Р	(N)	M1
Н	L	А	J	Ν	V	S	Т	Р	(L S)	A1
Н	А	J	L	N	Р	S	V	Т	(A V)	A1ft
А	Н	J	L	Ν	Р	S	Т	V	(H)	
А	Н	J	L	N	Р	S	Т	V		A1cso
First										
Н	V	L	А	Ν	J	S	Т	Р	(H)	M1
А	Н	V	L	Ν	J	S	Т	Р	(V)	A1
А	Н	L	Ν	J	S	Т	Р	V	(L)	
А	Н	J	L	Ν	S	Т	Р	V	(N)	Alft
A	Н	J	L	N	S	Т	Р	V	(S)	

[8]

5

4.	(a)										
0.6	4.0	2.5	3.2	0.5	2.6	0.4	0.3	4.0	1.0	2.6	
4.0	3.2	4.0	<u>2.6</u>	0.6	2.5	0.5	0.4	0.3	1.0	3.2 0.4	M1
4.0	4.0	<u>3.2</u>	<u>2.6</u>	0.6	2.5	0.5	1.0	<u>0.4</u>	0.3	4.0 0.5	A1
4.0	<u>4.0</u>	<u>3.2</u>	2.6	0.6	2.5	1.0	<u>0.5</u>	<u>0.4</u>	<u>0.3</u>	2.5	A1ft
4.0	<u>4.0</u>	<u>3.2</u>	2.6	<u>2.5</u>	0.6	1.0	<u>0.5</u>	<u>0.4</u>	<u>0.3</u>	1.0	A1ft
<u>4.0</u>	<u>4.0</u>	<u>3.2</u>	<u>2.6</u>	<u>2.5</u>	<u>1.0</u>	0.6	<u>0.5</u>	<u>0.4</u>	<u>0.3</u>		A1 cso

### <u>Notes</u>

	pivot per iteration M1 only
	<pre>p. If only choosing 1</pre>
1M1	Pivot, p, chosen. List sorted, >p, p.

- 1A1 1<sup>st</sup> pass correct and chosen next **two** pivots correctly for sublists >1
- 2A1ft 2<sup>nd</sup> pass correct and chosen next **two** pivots correctly for sublists >1
- 3A1ft  $3^{rd}$  pass correct and next pivot for sublist >1 chosen correctly.
- 4A1 cso.

### Misread in part (a)

- If they have misread a number **at the start of part (a), so genuinely miscopied** and got for example 0.1 instead of 1.0 then mark the whole question as a misread – removing the last two A or B marks earned. This gives a maximum total of 9.
- If they misread their own numbers during the course of part (a) then count it as an error in part (a) but mark parts (b) and (c) as a misread. So they would lose marks in (a) for the error and then the last two A or B marks earned in (b) and (c) – giving a maximum of 8 or maybe 7 marks depending on how many marks they lose in (a).

The most popular misread is the one listed above – where 1.0 has changed to 0.1 giving

4.0 4.0 3.2 2.6 2.5 0.6 0.5 0.4 0.3 <b>0.1</b>	at the end of (a) for
	this one (b) and (c)
	are:

(b)	Length 1:	4
(U)	Longui I.	-

Length 2:	4				
Length 3:	3.2	0.6	left colum	nn & 1.0 in place	M1
Length 4:	2.6	1.0	0.4	0.6 & 0.5	A1
Length 5:	2.5	0.5	0.3	0.4	A1

All correct (c.s.o)

A1

4

### <u>Note</u>

Length 1: 4 Length 2: 4 Length 3: 3.2 0.6 0.1 Length 4: 2.6 0.5 0.4 0.3 Length 5: 2.5

(c)	19.1/4 = 4.775 so 5 lengths needed, accept total is 19.1m, or refer to 0.9 'spare .	B1	
	Yes, the answer to (b) does use the minimum number of bins.	DB1	2
	Note		
	18.2/4 = 4.55 so 5 bins, or total is 18.2 or 1.8 'spare'		

Yes answer in (b) uses the minimum number of bins.

# Alternate

Choosi	ing mic	ldle lef	ť							
0.6	4.0	2.5	3.2	<u>0.5</u>	2.6	0.4	0.3	4.0	1.0	(pivot 0.5)
0.6	4.0	2.5	<u>3.2</u>	2.6	4.0	1.0	0.5	<u>0.4</u>	0.3	(pivots 3.2, 0.4)
4.0	4.0	3.2	0.6	<u>2.5</u>	2.6	1.0	0.5	0.4	<u>0.3</u>	(pivots 4.0, 2.5)
4.0	4.0	3.2	<u>2.6</u>	2.5	<u>0.6</u>	1.0	0.5	0.4	0.3	(pivot 0.6)
4.0	4.0	3.2	2.6	2.5	<u>1.0</u>	0.6	0.5	0.4	0.3	
4.0	4.0	3.2	2.6	2.5	1.0	0.6	0.5	0.4	0.3	
Choosi	ing firs	t								
<u>0.6</u>	4.0	2.5	3.2	0.5	2.6	0.4	0.3	4.0	1.0	(pivot 0.6)
4.0	2.5	3.2	2.6	4.0	1.0	0.6	<u>0.5</u>	0.4	0.3	(pivots 4.0, 0.5)
4.0	<u>2.5</u>	3.2	2.6	4.0	1.0	0.6	0.5	0.4	0.3	(pivots 2.5, 0.4)
4.0	<u>3.2</u>	2.6	4.0	2.5	1.0	0.6	0.5	0.4	<u>0.3</u>	(pivot 3.2)
4.0	<u>4.0</u>	3.2	<u>2.6</u>	2.5	1.0	0.6	0.5	0.4	0.3	
4.0	4.0	3.2	2.6	2.5	1.0	0.6	0.5	0.4	0.3	
OR (al	ternate	choosi	ing firs	t)						
<u>0.6</u>	4.0	2.5	3.2	0.5	2.6	0.4	0.3	4.0	1.0	(pivot 0.6)
$\frac{4.0}{1.0}$	2.5	3.2	2.6	4.0	1.0	0.6	<u>0.5</u>	0.4	0.3	(pivots 4.0, 0.5)
$\frac{4.0}{1.0}$	4.0	$\frac{2.5}{2.5}$	3.2	2.6	1.0	0.6	0.5	$\frac{0.4}{0.4}$	0.3	(pivots 2.5, 0.4)
4.0	4.0	<u>3.2</u>	2.6	2.5	$\frac{1.0}{1.0}$	0.6	0.5	0.4	$\frac{0.3}{0.2}$	(pivots 3.2)
4.0	4.0	3.2	$\frac{2.6}{2.6}$	2.5	1.0	0.6	0.5	0.4	0.3	
4.0	4.0	3.2	2.6	2.5	1.0	0.0	0.5	0.4	0.3	
Sorti	ng into	ASCE	ENDIN	G orde	er (full	marks	s if the	n reve	rsed, o	therwise MISREAD)
Middle	e left									
0.6	4.0	2.5	3.2	<u>0.5</u>	2.6	0.4	0.3	4.0	1.0	(pivot 0.5)
0.4	0.3	0.5	0.6	4.0	2.5	<u>3.2</u>	2.6	4.0	1.0	(pivots 0.4, 3.2)
<u>0.3</u>	0.4	0.5	0.6	<u>2.5</u>	2.6	1.0	3.2	<u>4.0</u>	<u>4.0</u>	(pivots 2.5, 4.0)

0.3

0.3

0.4

0.4

0.5

0.5

0.6

0.6

1.0

1.0

2.5

2.5

2.6

2.6

3.2

3.2

4.0

4.0

4.0

4.0

(pivot 0.6)

Middle	right									
0.6	4.0	2.5	3.2	0.5	2.6	0.4	0.3	4.0	1.0	(pivot 2.6)
0.6	2.5	0.5	<u>0.4</u>	0.3	1.0	2.6	4.0	<u>3.2</u>	4.0	(pivots 0.4, 3.2)
0.3	0.4	0.6	2.5	<u>0.5</u>	1.0	2.6	3.2	4.0	4.0	(pivots 0.5, 4.0)
0.3	0.4	0.5	0.6	<u>2.5</u>	1.0	2.6	3.2	4.0	4.0	(pivot 2.5)
0.3	0.4	0.5	0.6	1.0	2.5	2.6	3.2	4.0	4.0	(pivot 1.0)
First (1	)									
<u>0.6</u>	4.0	2.5	3.2	0.5	2.6	0.4	0.3	4.0	1.0	(pivot 0.6)
<u>0.5</u>	0.4	0.3	0.6	4.0	2.5	3.2	2.6	4.0	1.0	(pivots 0.5, 4.0)
<u>0.4</u>	0.3	0.5	0.6	2.5	3.2	2.6	1.0	4.0	4.0	(pivots 0.4, 2.5)
0.3	0.4	0.5	0.6	1.0	2.5	<u>3.2</u>	2.6	4.0	4.0	(pivot 3.2)
0.3	0.4	0.5	0.6	1.0	2.5	2.6	3.2	4.0	4.0	
First (2	)									
0.6	40	25	32	0.5	26	04	03	40	1.0	(nivot 0.6)
0.5	-1.0	0.3	0.6	4.0	2.0	3.2	2.6	4.0	1.0	(pivot 0.0)
$\frac{0.5}{0.4}$	0.4	0.5	0.6	$\frac{1.0}{2.5}$	$\frac{2.3}{3.2}$	2.6	$\frac{2.0}{1.0}$	4.0	4.0	(pivots 0.3, 4.0) (pivots 0.4, 2.5)
$\frac{0.1}{0.3}$	0.4	0.5	0.6	$\frac{2.5}{1.0}$	2.5	3.2	2.6	4.0	4.0	(pivot 3 0; 1, 2.0)
0.3	0.4	0.5	0.6	1.0	2.5	2.6	<b>3.2</b>	4.0	4.0	(P100 5.2)

# [11]

# **5.** (a)

				1			1			1	<b></b>
M1 1A	В	D	Ν	Р	L	В	Н	Κ	Е	J	Μ
	Н	D	Ν	Р	L	Н	K	E	J	М	B
2A1:	D L	Ν	Р	L	Κ	J	М	H	D	Е	B
КР	(E) K P	Ν	Р	М	L	K	J	H	E	D	B
N 3A1:	(J) N	P	Ν	М	L	K	J	H	E	D	B
	(M)	P	N	Μ	L	K	J	H	E	D	B
							-				

Sort completed

4A1 5

### <u>Note</u>

1M1: quick sort, pivots, p, identified, two sublists one p.

# If choosing one pivot only per iteration, M1 only.

- 1A1: first pass correct, next pivot(s) chosen consistently.
- 2A1ft: second pass correct, next pivot(s) chosen consistently
- 3A1ft: third pass correct, next pivot(s) chosen consistently
- 4A1: cso List re-written or end statement made or each element been chosen as a pivot.

(b) 
$$\left[\frac{1+10}{2}\right] = 6$$
 Katie reject left M1  
 $\left[\frac{7+10}{2}\right] = 9$  Natsuko reject right 1A1  
 $\left[\frac{7+8}{2}\right] = 8$  Miri reject right 2A1ft  
7 = Louis name found 3A1 4

### <u>Note</u>

1M1: binary search, choosing pivot rejecting half list.

### If using unordered list then M0.

### If choosing J M1 ony

- 1A1: first two passes correct, condone 'sticky'pivots here, bod.
- 2A1ft: third pass correct, pivots rejected.
- 3A1: cso, including success statement.

### Special case

If just one letter out of order, award maximum of M1A1A0A0

[9]

### **6.** (a) e.g.



A1

4

Η	J	Ι	K	М	L	R	Т	A1ft	
Η	Ι	J	Κ	L	М	R	Т	A1ft	
Η	Ι	J	K	L	М	R	Т	A1cso	5

### <u>Note</u>

1M1: quick sort, pivots, p, chosen and two sublists one p. If choosing 1 pivot per iteration only M1 only.

1A1: first pass correct and next pivots chosen correctly/consistently.2A1ft: second pass correct, next pivots correctly/consistently chosen.3A1ft: third pass correct, next pivots correctly/consistently chosen.4A1: all correct, cso.

(b)

Sort complete.

$$1^{st} \text{ choice } \left[\frac{1+8}{2}\right] \rightarrow 5 \text{ Lauren reject right}$$
M1 A1  

$$2^{nd} \text{ choice } \left[\frac{1+4}{2}\right] \rightarrow 3 \text{ John reject right}$$
M1 A1  

$$3^{rd} \text{ choice } \left[\frac{1+2}{2}\right] \rightarrow 2 \text{ Imogen reject right}$$
A1ft

4<sup>th</sup> choice 1 Hannah reject List now empty so Hugo not in list

### <u>Note</u>

1M1: binary search, choosing pivot, rejecting half list. If using unsorted list, M0. Accept choice of K for M1 only.
1A1: first pass correct, condone 'sticky' pivot here, bod.
2A1ft: second pass correct, pivot rejected.
3A1: cso.

[9]

7. 
$$\left[\frac{1+10}{2}\right] = 6$$
 Nicky – reject top of list. M1  
 $\left[\frac{7+10}{2}\right] = 9$  Trevor – reject bottom of list A1  
 $\left[\frac{7+8}{2}\right] = 8$  Steve – reject bottom of list A1  
[7] = 7 Preety – reject A1  
Nigel not in list A1

8.	(a)	eg.	45	37	18	46	56	79	90	81	51
		or	37	18	45	56	79	46	90	81	51

[4]

	or	45	37	46	18	51	56	79	90	81	M1A1	2
(b)	or	56 90	45 45	79 56	46 37	37 79	90 46	81 18	51 81	18 51	M1A1	2
(c)	$\left[\frac{1+1}{2}\right]$	$\left\lfloor 1 \right\rfloor = 6$	value	44 dis	card to	р					M1	
	$\left[\frac{7+}{2}\right]$	$\left[\frac{11}{2}\right] = 9$	9 value	71 dis	scard to	op					A1	
	$\left[\frac{10+2}{2}\right]$	$\left[\frac{-11}{2}\right] =$	11 val	ue 94 o	discard	botto	m				A1	
	list re <u>73 ha</u>	educes <u>s been</u>	to 10 <sup>th</sup> locate	value. d as th	This <u>i</u> e 10 <sup>th</sup>	<u>s</u> 73 so value	)				A1	4

9.	(a)	The list is not in <u>alphabetical</u> order	B1
	(b)	Use of Bubble Sort or Quick Sort	M1

(b) Use of Bubble Sort or Quick Sort e.g.

### Bubble sort

_							_	-		
G	Ν	Μ	Y	L	В	C	E	S	Р	
В	G	Ν	Μ	Y	L	С	E	P	S	1 <sup>st</sup> pass
В	C	G	Ν	M	Y	L	E	P	S	2 <sup>nd</sup> pass
В	C	E	G	N	Μ	Y	L	P	S	3 <sup>rd</sup> pass
В	С	Е	G	L	Ν	Μ	Y	P	S	4 <sup>th</sup> pass
В	С	Е	G	L	Μ	Ν	Р	Y	S	5 <sup>th</sup> pass
В	C	E	G	L	Μ	Ν	Р	S	Y	6 <sup>th</sup> pass

No more changes

## Quick sort

G	Ν	Μ	Y	L	B	C	E	S	Р	
В	G	Ν	Μ	Y		C	E	S	Р	1 <sup>st</sup> pass
В	G	$\bigcirc$	E	L	Ν	Μ	(Y)	S	Р	2 <sup>nd</sup> pass
В	С	G	Œ	L	Ν	Μ	$(\mathbb{S})$	Р	Y	3 <sup>rd</sup> pass
В	С	Е	G	L	Ν	$\mathbb{M}$	Р	S	Y	4 <sup>th</sup> pass
В	С	Е	G	L	Μ	Ν	$\bigcirc$	S	Y	5 <sup>th</sup> pass
В	С	Е	G	L	Μ	Ν	Р	S	Y	6 <sup>th</sup> pass

No sublists > 2 and no more changes

No more changes

No sublists > 2 + no more changes

A1 A1ft A1cso

4

[8]

1

(c)	1	2	3	4	5	6	7	8	9	10			
	В	С	E	G	L	М	Ν	Р	S	Y			
	<u>[10 -</u> 2	+1] =	6	Manche	disc	ard firs	st half	M1 A1					
	$\frac{[7+10]}{2} = 9$		9	Southar	disc	ard las	t half (						
	$\frac{[7+2]{2}}{2}$	$\frac{8]}{8} = 8$	3	Plymou	th	disc	ard las	t half (	of list a	and pivot	A1ft		
	Final term 7 Newcastle ∴ word found at 7										A1cso	4	

### **10.** (a) e.g.

5

(b)  $\left[\frac{10+1}{2}\right] = 6$  Palmer; reject Palmer  $\rightarrow$  Young M1 A1  $\left[\frac{5+1}{2}\right] = 3$  Halliwell; reject Boase  $\rightarrow$  Halliwell A1  $\left[\frac{4+5}{2}\right] = 5$  Morris; reject Morris List reduces to Kenney – name found, search complete A1 4

[9]

[9]

- 1. No Report available for this question.
- 2. No Report available for this question.
- 3. This proved a good starter and was well answered by many candidates with around 55% getting full marks. The quick sort was well handled although some candidates did not choose their pivots consistently. A few candidates did not select a pivot when they had a two element sublist in the correct order often HJ, and a minority sorted the list into reverse alphabetical order. It was alarming that some candidates only selected one pivot per iteration, so, in effect, just dealing with one sublist at a time. Candidates must show that they are selecting one pivot, per sublist, per iteration; that is what makes this algorithm so powerful. A number of candidates did not have the final list in alphabetical order.

Many candidates in part (b) lost marks for failing to reject the pivot and number of candidates attempted to use the original, unsorted list. Some, who tried for a more 'minimalist' solution, did not make their pivot choice clear, or the order in which they chose pivots.

- 4. Many candidates scored at least 8 marks here. In part (a) a minority produced an ascending list and failed to reverse it. Some candidates did not choose their pivots consistently, swapping between middle right and middle left pivots. The decimals here caused some problems and even though the original list was printed in the answer booklet, a surprising number of candidates initially lost one item or changed one, most commonly 1.0 became 0.1. Some candidates found only one pivot per row, with some not explicitly choosing pivots when sublists of length 2 happened to be in order most frequently the two 4.0s and the 1.0, 0.6 at the end. Good presentation, with a list spread evenly, in columns, across the page, helps here. (Vertical listing is rarely successful). Part (b) was generally well done, the two most popular errors being to put 0.6 in bin 5 or 0.4 in bin 5. A significant number who had sorted the numbers into increasing order in part (a) proceeded to use a "first fit increasing" method here. In part (c) most candidates calculated the lower bound correctly. Other candidates correctly stated that since the five largest items were over half a bin in size they could not share a bin, so at least 5 bins would be needed. A few simply stated 'yes' without justification, gaining no credit.
- 5. This was generally well done. A disappointingly large number of candidates only chose one pivot per iteration, rather than choosing one pivot per sublist, and some candidates used lengthy methods of presentation that isolated each sublist in turn, making it difficult to see if they were choosing more than one pivot per iteration. The examiners would advise candidates to refrain from showing this unnecessary detail and simply indicate the pivots selected at each iteration. Some candidates did not select a pivot where the sublist was of order two, with the two items being in the correct order, and some did not consistently pick 'middle left' or 'middle right' when the sublist was of even order. Candidates are reminded that when the items are being transferred to the next line, the order of the items should be preserved, so if item Y is to the left of item X in the current line, neither of them being a pivot, then Y should be to the left of X in the next line. The best candidates allowed each item to become a pivot before declaring the sort complete. Some candidates did not check that their final list was in alphabetical order. In part (b) some candidates tried to apply the algorithm to the original unsorted list given at the start of (a) and others did not discard the pivot at each stage, but generally the binary search was very well done. A few candidates selected J as the first pivot, the specification makes it clear that candidates must take the 'middle right' where necessary.

6. Part (a) was done with mixed success. The majority of candidates gained full marks or three marks. The most common errors were to have HIJ after the second pass and neglecting to choose a pivot on the third pass with the entry MR. Most knew their alphabet, but not all. There was a temptation to go into too much detail about the choice of pivot, to the extent that examiners were not always sure that more than one pivot was being considered per iteration. It is an important feature of the quick sort that the number of pivots can potentially double at each iteration, so the selection of multiple pivots must be clearly shown. Some candidates did not abbreviate the names, by using the initial letter and this slowed them down.

Part (b) was usually very well done. The most common errors were not rejecting the pivot and not making a decision when Hannah was left. Some candidates added Hugo to the list and then found him, others confused Hannah and Hugo.

- 7. This proved a good starter question for the candidates with many gaining full marks. Some candidates were inconsistent in their pivot choice, the specification requires that they round up. Some incorrectly retained the pivot each time often leading to a situation where they selected Nicky twice, once as the first pivot and once as the final pivot. Some candidates insisted on placing Nigel in the list or locating the position in which Nigel should be added to the list. The binary search algorithm is both used to locate an item in the list and to demonstrate its absence. A few candidates confused binary search and quick sort.
- 8. This was generally well done. Many candidates completed the quick sort, wasting time. Some candidates did not understand the difference between an exchange and a pass in a bubble sort. Most candidates carried out the search well, but many did not give the location of the value. A large number are still assuming that the item is in the list, making statements such as 'down to one item so found'. A surprisingly large minority of candidates used the mean of the end numbers in the remaining list to create a 'pivot' which is unacceptable.

- 9. This question was often well answered. Most candidates correctly competed part (a), although a very few stated that the list should be in ascending rather than alphabetical order. Most could correctly name and use a suitable sorting algorithm in part (b), although some did not make their stopping statement clear and a few used a shuttle sort (not in this specification) stating that it was a bubble sort. A surprisingly large minority confused the order of the alphabet with S and P (and then M and N) most frequently transposed. Part (c) was usually well done but candidates must make their pivots and the order in which they select their pivots, clear. Candidates must remember to discard their pivots and note that the specification instructs them to 'round up'. Once again the stopping/found statement was sometimes missing, and some candidates assumed the presence of N, stating that once they had got down to 1 term only, that term must be N.
- 10. Many candidates were able to gain full marks on this question. The most common errors in part (a) were in re-ordering the letters in the sub-lists and choosing the pivots inconsistently. A surprising number of candidates seemed unsure of the alphabet. Part (b) was well done by the majority of candidates. A surprising number tried to use an unsorted list for their search, gaining no marks and others omitted to discard the pivot. The commonest error was in failing to select Morris after correctly selecting Palmer then Halliwell. A few candidates did not make the order in which they selected the pivots clear making it impossible to give credit.